

## RAA 2017

# Longest Paths in Graphs: Parameterized Algorithms – Exercise I

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**Remark :** Several of these problems use the notion of *matroids*. Those who do not know the definition of matroids can find the definitions below. Else, assume that  $\mathcal{S}$  is a family of sets over a universe  $U$  such that each set in  $\mathcal{S}$  has size exactly  $p$ .

1. In the VERTEX COVER problem, we are given a graph  $G = (V, E)$  and a positive integer  $k$ , and the problem is to test whether there exists a vertex subset  $X \subseteq V(G)$  such that  $|X| \leq k$  and  $G \setminus X$  is an independent set. Obtain a  $2^k n^{\mathcal{O}(1)}$ -time algorithm for this problem.
2. In the  $d$ -HITTING SET problem, we are given a family  $\mathcal{F}$  of sets of size  $d$  over a universe  $U$  and a positive integer  $k$ , and the problem is to test whether there exists a subset  $X \subseteq U$  such that  $|X| \leq k$  and for every set  $F \in \mathcal{F}$ ,  $F \cap X \neq \emptyset$ . Obtain a  $k^{\mathcal{O}(d)}$  kernel for the problem using the method of representative sets.
3. Let  $A_1, \dots, A_m$  be  $p$ -element sets and  $B_1, \dots, B_m$  be  $q$ -element sets such that  $A_i \cap B_j = \emptyset$  if and only of  $i = j$ .
  - (a) Show that  $m \leq 2^{p+q}$ . (Hint: Think uniform random partition of  $U = \cup_{i=1}^m (A_i \cup B_i)$ .)
  - (b) Show that  $m \leq \binom{p+q}{p}$ . (Hint: Think of permutations of  $U$ .)
  - (c) Show that the bound of  $\binom{p+q}{p}$  on  $m$  is tight.
  - (d) Let  $\mathcal{S} = \{S_1, \dots, S_t\}$  be a family of  $p$  element sets. Using the above exercises show that the size of  $q$ -representative family is upper bounded by  $\binom{p+q}{p}$ .
4. Let  $M = (U, \mathcal{I})$  be a matroid and let  $\mathcal{S}$  be a  $p$ -uniform family of subsets of  $E$ . Show that if  $\mathcal{S}' \subseteq_{rep}^q \mathcal{S}$  and  $\widehat{\mathcal{S}} \subseteq_{rep}^q \mathcal{S}'$ , then  $\widehat{\mathcal{S}} \subseteq_{rep}^q \mathcal{S}$ . (If  $\widehat{\mathcal{S}} \subseteq \mathcal{S}$  is  $q$ -representative for  $\mathcal{S}$  we write  $\widehat{\mathcal{S}} \subseteq_{rep}^q \mathcal{S}$ .)
5. Let  $M = (U, \mathcal{I})$  be a matroid and let  $\mathcal{S}$  be a  $p$ -uniform family of subsets of  $E$ . Show that if  $\mathcal{S} = \mathcal{S}_1 \cup \dots \cup \mathcal{S}_\ell$  and  $\widehat{\mathcal{S}}_i \subseteq_{rep}^q \mathcal{S}_i$ , then  $\cup_{i=1}^\ell \widehat{\mathcal{S}}_i \subseteq_{rep}^q \mathcal{S}$ .
6. Let  $G$  be a connected graph on  $2n$  vertices and  $\mathcal{L}$  be a family of forests of  $G$  of size  $n$  (that is, the number of edges is  $n$ ). Let  $\widehat{\mathcal{L}} \subseteq \mathcal{L}$  be a family of forests such that for any forest  $F$  of size  $n - 1$ , if there exists a forest  $X \in \mathcal{L}$  such that  $F \cup X$  is a spanning tree of  $G$ , then there exists a forest  $\widehat{X} \in \widehat{\mathcal{L}}$  such that  $F \cup \widehat{X}$  is a spanning tree of  $G$ . Could you give a non-trivial upper bound on the size of  $|\widehat{\mathcal{L}}|$  (like some  $c^n$ )?

## Matroid Basics

Now we give definitions related to matroids.

**Definition 1** A pair  $M = (U, \mathcal{I})$ , where  $E$  is a ground set and  $\mathcal{I}$  is a family of subsets (called independent sets) of  $E$ , is a matroid if it satisfies the following conditions:

- (I1)  $\phi \in \mathcal{I}$ .
- (I2) If  $A' \subseteq A$  and  $A \in \mathcal{I}$  then  $A' \in \mathcal{I}$ .
- (I3) If  $A, B \in \mathcal{I}$  and  $|A| < |B|$ , then there is  $e \in (B \setminus A)$  such that  $A \cup \{e\} \in \mathcal{I}$ .

The axiom (I2) is also called the hereditary property and a pair  $(E, \mathcal{I})$  satisfying only (I2) is called hereditary family. An inclusion wise maximal set of  $\mathcal{I}$  is called a *basis* of the matroid. Using axiom (I3) it is easy to show that all the bases of a matroid have the same size. This size is called the *rank* of the matroid  $M$ , and is denoted by  $\text{rank}(M)$ .

## Linear Matroids and Representable Matroids

Let  $A$  be a matrix over an arbitrary field  $\mathbb{F}$  and let  $E$  be the set of columns of  $A$ . For  $A$ , we define matroid  $M = (U, \mathcal{I})$  as follows. A set  $X \subseteq E$  is independent (that is  $X \in \mathcal{I}$ ) if the corresponding columns are linearly independent over  $\mathbb{F}$ . The matroids that can be defined by such a construction are called *linear matroids*, and if a matroid can be defined by a matrix  $A$  over a field  $\mathbb{F}$ , then we say that the matroid is representable over  $\mathbb{F}$ . That is, a matroid  $M = (U, \mathcal{I})$  of rank  $d$  is representable over a field  $\mathbb{F}$  if there exist vectors in  $\mathbb{F}^d$  corresponding to the elements such that linearly independent sets of vectors correspond to independent sets of the matroid. A matroid  $M = (U, \mathcal{I})$  is called *representable* or *linear* if it is representable over some field  $\mathbb{F}$ .

1. Show that the following families form matroid.
  - (a) Let  $G = (V, E)$  be a graph. Let  $M = (U, \mathcal{I})$  be a matroid defined on  $G$ , where  $U = E$  and  $\mathcal{I}$  contains all *forests* of  $G$ . (**Graphic Matroid**)
  - (b) Let  $G = (V, E)$  be a connected graph. Let  $M = (U, \mathcal{I})$  be a matroid defined on  $G$ , where  $U = E$  and  $\mathcal{I}$  contains all  $E' \subseteq E$  such that  $G' = (V, E \setminus E')$  is connected. (**Co-Graphic Matroid**)
2. Obtain a representation matrix for the following matroid.
  - (a) Graphic Matroid.
  - (b) Uniform Matroids –  $M = (U, \mathcal{I})$  where  $\mathcal{I}$  contains all subsets of  $U$  of size at most  $k$  for some fixed constant  $k$ .
  - (c) Partition Matroids – It is defined by a ground set  $U$  being partitioned into (disjoint) sets  $U_1, \dots, U_\ell$  and by  $\ell$  non-negative integers  $k_1, \dots, k_\ell$ . A set  $X \subseteq U$  is independent if and only if  $|X \cap U_i| \leq k_i$  for all  $i \in \{1, \dots, \ell\}$ . That is,

$$\mathcal{I} = \left\{ X \subseteq U \mid |X \cap U_i| \leq k_i, i \in \{1, \dots, \ell\} \right\}.$$

- (d) Direct Sum of Matroids – Let  $M_1 = (U_1, \mathcal{I}_1)$ ,  $M_2 = (U_2, \mathcal{I}_2)$ ,  $\dots$ ,  $M_t = (U_t, \mathcal{I}_t)$  be  $t$  matroids with  $U_i \cap U_j = \emptyset$  for all  $1 \leq i \neq j \leq t$ . The direct sum  $M_1 \oplus \dots \oplus M_t$  is a matroid  $M = (U, \mathcal{I})$  with  $U := \bigcup_{i=1}^t U_i$  and  $X \subseteq U$  is independent if and only if for all  $i \leq t$ ,  $X \cap U_i \in \mathcal{I}_i$ .
3. Let  $M_1 = (U_1, \mathcal{I}_1)$  and  $M_2 = (U_2, \mathcal{I}_2)$  be two matroids such that  $U = U_1 = U_2$ . Define  $M_1 \cap M_2$  as  $M = (U, \mathcal{I})$  such that  $X \in \mathcal{I}$  if and only if  $X \in \mathcal{I}_1$  and  $X \in \mathcal{I}_2$ . Is  $M$  always a matroid? (**Matroid Intersection**)
4. Express the following as intersection of matroids (possibly more than two).
- (a) Finding a maximum matching in a bipartite graph  $G = (A \cup B, E)$ .
  - (b) Testing whether a graph  $G = (V, E)$  contains two edge disjoint spanning trees.
  - (c) Finding a hamiltonian path in a directed graph  $D = (V, A)$  between a pair of vertices  $s$  and  $t$  of  $D$ .